

## A. Appendix: Trees and Majorizing Measures

In this appendix we describe different ways to measure the size of a metric space. Some of these ways played an important part in the development of the theory. We will show that they are all equivalent to the functional  $\gamma_2(T, d)$ . It is possible to consider more general notions corresponding to other functionals considered in the book, but for simplicity we consider only the case of  $\gamma_2$ .

A *tree*  $\mathcal{T}$  of a metric space  $(T, d)$  is a *finite* collection of subsets of  $T$  with the following two properties.

Given  $A, B$  in  $\mathcal{T}$ , if  $A \cap B \neq \emptyset$ , then either  $A \subset B$  or else  $B \subset A$ . (A.1)

$\mathcal{T}$  has a largest element. (A.2)

If  $A, B \in \mathcal{T}$  and  $B \subset A$ ,  $B \neq A$ , we say that  $B$  is a *child* of  $A$  if

$C \in \mathcal{T}$ ,  $B \subset C \subset A \Rightarrow C = B$  or  $C = A$ . (A.3)

We denote by  $c(A)$  the number of children of  $A$ . We will consider only trees with the following property

If  $A \in \mathcal{T}$  and  $c(A) = 0$ , then  $A$  contains exactly one point. (A.4)

A *separated* tree is a tree  $\mathcal{T}$  such that to each  $A$  in  $\mathcal{T}$  with  $c(A) \geq 1$  is associated an integer  $s(A) \in \mathbb{Z}$  with the following properties.

If  $B_1$  and  $B_2$  are children of  $A$ , then  $d(B_1, B_2) \geq 4^{-s(A)}$ . (A.5)

If  $B$  is a child of  $A$ , then  $s(B) > s(A)$ . (A.6)

Here of course  $d(B_1, B_2) = \inf\{d(x_1, x_2), x_1 \in B_1, x_2 \in B_2\}$ . We then define

$$S_{\mathcal{T}} = \{t \in T ; \{t\} \in \mathcal{T}\}$$

and the *depth* of  $\mathcal{T}$ ,

$$d(\mathcal{T}) = \inf_{t \in S_{\mathcal{T}}} \sum_{t \in A \in \mathcal{T}} 4^{-s(A)} \sqrt{\log c(A)}.$$

Here and below we make the convention that the summation does not include the term  $A = \{t\}$  (for which  $c(A) = 0$ ). We view  $d(\mathcal{T})$  as a “measure

of the size” of the separated tree  $\mathcal{T}$ . We can then measure the size of  $T$  by  $\sup\{d(\mathcal{T}); \mathcal{T} \text{ separated tree}\}$ .

The notion of tree we just considered is but one of many possible. Let us now consider another (more restrictive) notion. An *organized* tree is a tree  $\mathcal{T}$  such that to each  $A \in \mathcal{T}$  with  $c(A) \geq 1$  are associated  $j = j(A) \in \mathbb{Z}$ ,  $t \in T$  and  $t_1, \dots, t_{c(A)} \in B(t, 4^{-j})$  with the properties that

$$1 \leq \ell < \ell' \leq c(A) \Rightarrow d(t_\ell, t_{\ell'}) \geq 4^{-j-1}$$

and that each ball  $B(t_\ell, 4^{-j-2})$  contains exactly one child of  $A$ .

If  $B_1$  and  $B_2$  are children of  $A$ , then

$$d(B_1, B_2) \geq 4^{-j(A)-2}, \quad (\text{A.7})$$

so that an organized tree is also a separated tree (with  $s(A) = j(A) + 2$ ), but the notion of organized tree is more restrictive. (For example we have no control of the diameter of the children of  $A$  in a separated tree.)

We define the depth  $d'(\mathcal{T})$  of an organized tree by

$$d'(\mathcal{T}) = \inf_{t \in S_{\mathcal{T}}} \sum_{t \in A \in \mathcal{T}} 4^{-j(A)} \sqrt{\log c(A)}.$$

If we simply view  $\mathcal{T}$  as a separated tree using (A.7), then  $d(\mathcal{T}) = d'(\mathcal{T})/16$  (where  $d(\mathcal{T})$  is the depth of  $\mathcal{T}$  as a separated tree). Thus we have shown the following.

**Proposition A.1.** *We have*

$$\sup\{d'(\mathcal{T}); \mathcal{T} \text{ organized tree}\} \leq 16 \sup\{d(\mathcal{T}); \mathcal{T} \text{ separated tree}\}. \quad (\text{A.8})$$

**Proposition A.2.** *We have*

$$\gamma_2(T, d) \leq L \sup\{d'(\mathcal{T}); \mathcal{T} \text{ organized tree}\}. \quad (\text{A.9})$$

*Proof.* We consider the functional

$$F_n(A) = F(A) = \sup\{d'(\mathcal{T}); \mathcal{T} \subset A, \mathcal{T} \text{ organized tree}\},$$

where we write  $\mathcal{T} \subset A$  as a shorthand for “ $\forall B \in \mathcal{T}, B \subset A$ ”. In the course of the proof of Theorem 1.3.1 we have noted that this theorem holds true as soon as (1.31) holds true when  $a$  is of the type  $r^{-j-1}$ . We check this condition when  $r = 4$ ,  $\theta(n) = 2^{n/2-2}$ ,  $\beta = 1$ , and  $\tau = 1$ . Consider  $n \geq 0$  and  $m = N_{n+1}$ . Consider  $j \in \mathbb{Z}$ ,  $t \in T$  and  $t_1, \dots, t_m \in B(t, 4^{-j})$  with

$$1 \leq \ell < \ell' \leq m \Rightarrow d(t_\ell, t_{\ell'}) \geq 4^{-j-1}.$$

Consider sets  $H_\ell \subset B(t_\ell, 4^{-j-2})$  and  $c < \min_{\ell \leq m} F(H_\ell)$ . Consider, for  $\ell \leq m$  a tree  $\mathcal{T}_\ell \subset H_\ell$  with  $d'(\mathcal{T}_\ell) > c$  and denote by  $A_\ell$  its largest element. Then it should be obvious that the tree  $\mathcal{T}$  consisting of  $U = \bigcup_{\ell \leq m} H_\ell$  (its

largest element) and the reunion of the trees  $\mathcal{T}_\ell$ ,  $\ell \leq m$ , is organized (with  $j(U) = j$ , and  $A_1, \dots, A_m$  as children of  $U$ ). Moreover  $S_{\mathcal{T}} = \bigcup_{\ell \leq m} S_{\mathcal{T}_\ell}$ .

Consider  $t \in S_{\mathcal{T}}$ , and let  $\ell$  with  $t \in S_{\mathcal{T}_\ell}$ . Then

$$\begin{aligned} \sum_{t \in A \in \mathcal{T}} 4^{-j(A)} \sqrt{\log c(A)} &= 4^{-j} \sqrt{\log m} + \sum_{t \in A \in \mathcal{T}_\ell} 4^{-j(A)} \sqrt{\log c(A)} \\ &\geq 4^{-j} \sqrt{\log m} + d'(\mathcal{T}_\ell) \geq 4^{-j} \sqrt{\log m} + c. \end{aligned}$$

Since  $\sqrt{\log m} \geq 2^{n/2}$ , this proves (1.31).

To prove (A.9) we apply Lemma 1.3.3. To control the diameter of  $T$ , we simply note that if  $s, t \in T$ , and  $j$  is the largest integer with  $4^{-j} \geq d(s, t)$ , then the tree  $\mathcal{T}$  consisting of  $T, \{t\}, \{s\}$ , is organized with  $j = j(T)$  and  $c(T) = 2$ , so  $d'(\mathcal{T}) \geq 4^{-j} \sqrt{\log 2}$ .  $\square$

**Proposition A.3.** *Given a metric space  $(T, d)$  we can find on  $T$  a probability measure  $\mu$ , supported by a countable subset of  $T$ , and such that*

$$\sup_{t \in T} \int_0^\infty \sqrt{\frac{1}{\log \mu(B(t, \epsilon))}} d\epsilon \leq L\gamma_2(T, d). \quad (\text{A.10})$$

A probability measure  $\mu$  on  $(T, d)$  such that the left-hand side of (A.10) is usefully small is called a majorizing measure. The (in)famous theory of majorizing measures used the infimum of the left-hand side of (A.10) over all choices of  $\mu$  as a measure of the size of the metric space  $(T, d)$ .

*Proof.* Consider an admissible sequence  $(\mathcal{A}_n)$  with

$$\forall t \in T, \sum_{n \geq 0} 2^{n/2} \Delta(\mathcal{A}_n(t)) \leq 2\gamma_2(T, d).$$

Since  $\text{card } \mathcal{A}_n \leq N_n$ , there is a probability measure  $\mu$  on  $T$ , supported by a countable set, and satisfying

$$\forall n \geq 1, \forall A \in \mathcal{A}_n, \mu(A) \geq \frac{1}{2^n N_n} \geq \frac{1}{N_n^2}$$

so that given  $t \in T$

$$\begin{aligned} \epsilon > \Delta(\mathcal{A}_n(t), d) &\Rightarrow \mu(B(t, \epsilon)) \geq \frac{1}{N_n^2} \\ &\Rightarrow \sqrt{\log \frac{1}{\mu(B(t, \epsilon))}} \leq 2^{n/2+1}. \end{aligned} \quad (\text{A.11})$$

Now, since  $\mu$  is a probability,  $\mu(B(t, \epsilon)) = 1$  for  $\epsilon > \Delta(T, d)$ , and thus  $\log(1/\mu(B(t, \epsilon))) = 0$ . Thus

$$\begin{aligned} \int_0^\infty \sqrt{\log \frac{1}{\mu(B(t, \epsilon))}} d\epsilon &= \sum_{n \geq 1} \int_{\Delta(A_n(t))}^{\Delta(A_{n-1}(t))} \sqrt{\log \frac{1}{\mu(B(t, \epsilon))}} d\epsilon \\ &\leq \sum_{n \geq 1} 2^{n/2+1} \Delta(A_{n-1}(t)) \leq L\gamma_2(T, d) \end{aligned}$$

using (A.11).  $\square$

**Proposition A.4.** *If  $\mu$  is a probability measure on  $T$ , (supported by a countable set) and  $\mathcal{T}$  is a separated tree on  $T$ , then*

$$d(\mathcal{T}) \leq L \sup_{t \in T} \int_0^\infty \sqrt{\log \frac{1}{\mu(B(t, \epsilon))}} d\epsilon .$$

This completes the proof that the 4 “measures of the size of  $T$ ” considered in this appendix are indeed equivalent.

*Proof.* The basic observation is as follows. The sets

$$B(C, 4^{-s(A)-1}) = \{x \in T ; d(x, C) \leq 4^{-s(A)-1}\}$$

are disjoint as  $C$  varies over the children of  $A$  (as follows from (A.5)). So one of them has measure  $\leq c(A)^{-1}$ .

We then proceed in the following manner. We start with the largest element  $A_0$  of  $\mathcal{T}$ . We then select a child  $A_1$  of  $A_0$  with  $\mu(B(A_1, 4^{-s(A_0)-1})) \leq 1/c(A_0)$ , and a child  $A_2$  of  $A_1$  with  $\mu(B(A_2, 4^{-s(A_1)-1})) \leq 1/c(A_1)$ , etc., and continue this construction as long as we can. It ends only when we reach a set of  $\mathcal{T}$  that has no child, and hence by (A.4) is reduced to a single point  $t$ . If  $t \in A \in \mathcal{T}$ , by construction we have

$$\mu(B(t, 4^{-s(A)-1})) \leq \frac{1}{c(A)}$$

so that

$$4^{-s(A)-2} \sqrt{\log c(A)} \leq \int_{4^{-s(A)-2}}^{4^{-s(A)-1}} \sqrt{\log \frac{1}{\mu(B(t, \epsilon))}} d\epsilon . \quad (\text{A.12})$$

By (A.6) the intervals  $]4^{-s(A)-2}, 4^{-s(A)-1}[$  are disjoint for different sets  $A$  with  $t \in A \in \mathcal{T}$ , so summation of the inequalities (A.12) yields

$$\frac{1}{16} d(\mathcal{T}) \leq \sum_{t \in A \in \mathcal{F}} 4^{-s(A)-2} \sqrt{\log c(A)} \leq \int_0^\infty \sqrt{\log \frac{1}{\mu(B(t, \epsilon))}} d\epsilon .$$

$\square$